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CS193J: Programming in Java
Summer Quarter 2003

Lecture 10
Thread Interruption, Cooperation (wait/notify),
Swing Thread, Threading conclusions

Manu Kumar
sneaker@stanford.edu


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Handouts

- 1 Handout for today!
 - #21: Threading 3
 - #22: HW3a: ThreadBank
 - #23: HW3b: LinkTester

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Homework #2 feedback

- What did you think?
 - SCPD students are again encouraged to email their comments to me at sneaker@stanford.edu


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Recap

- Last Time
 - Review Introduction to Threading
 - Java threads
 - Simple Thread Example
 - Threading 2
 - Race Conditions
 - Reader/Writer Conflict
 - Writer/Writer Conflict
 - Locking
 - Synchronized Method
 - Synchronized method example


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Today

- Homework #3 overview
 - ThreadBank
 - demo
 - LinkTester
 - demo
- Thread Interruption
- Cooperation
 - Wait/notify
- Swing Thread
- Threading conclusions

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HW3a: ThreadBank

- Small assignment
 - Intended to cover mostly material covered in lecture this week
 - Java Threads
 - Synchronization
 - Cooperation (today)
- Recommendation
 - Finish this assignment this week!
 - It is small, the material is fresh in your mind
 - Part 3b is more involved...

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HW3b: LinkTester

- Based on the following material
 - Threading
 - Basic Networking
 - So basic that we will not cover this in lecture in detail – just a simple example
 - See handout and refer to API classes
 - Model-View-Controller
 - Swing Tables
 - Swing Thread
- Demo of HW3b...



Thread Interruption

- interrupt()
 - Signal a thread object that it should stop running
 - Asynchronous notification
 - Does not stop the thread right away
 - Sets an "interrupted" boolean to true
 - Thread must check and do appropriate thing
- isInterrupted()
 - Checks to see if a interrupt has been requested
 - Idiom – check isInterrupted() in a loop
 - When interrupted, should exit leaving object in a clean state



Stop() -- deprecated

- stop()
 - Performs a synchronous stop of the thread
 - Usually impossible to ensure that the object is left in a consistent state when using stop
 - Deprecated in favor of using interrupt() and doing a graceful exit



Interruption() example

```
class StopWorker extends Thread {
    public void run() {
        long sum = 0;
        for (int i=0; i<5000000; i++) {
            sum = sum + i; // do some work
            // every n iterations... check isInterrupted()
            if (i%100000 == 0) {
                if (isInterrupted()) {
                    // clean up, exit when interrupted
                    // (getName() returns a default name for each thread)
                    System.out.println(getName() + " interrupted");
                    return;
                }
                System.out.println(getName() + " " + i);
                Thread.yield();
            }
        }
    }
}
```



Interruption() example

```
public static void main(String[] args) {
    StopWorker a = new StopWorker();
    StopWorker b = new StopWorker();

    System.out.println("Starting...");
    a.start();
    b.start();
    try {
        Thread.sleep(100); // sleep a little, so they make some progress
    } catch (InterruptedException ignored) {}

    a.interrupt();
    b.interrupt();
    System.out.println("Interruption sent");
    try {
        a.join();
        b.join();
    } catch (Exception ignored) {}
    System.out.println("All done");
}
```



Interruption() example output

```
• /*
• Starting...
• Thread-0 0
• Thread-1 0
• Thread-1 100000
• Thread-0 100000
• Thread-1 200000
• ...
• Thread-0 900000
• Interruption sent
• Thread-0 interrupted
• Thread-1 interrupted
• All done
• */
```



- Threading Challenges
 - Synchronization
 - Preventing threads from stepping on each other when dealing with shared memory
 - Done using synchronized methods and synchronized(obj) {...} constructs
 - Cooperation/Coordination
 - Making on thread wait for the other
 - Signaling between threads
 - Done using join(), wait() and notify() constructs
 - join() we have already seen.



- Suppose we want to execute the statement
 - if (len > 0) len ++
- Problems:
 - Multiple threads
 - The statement is not atomic
 - The value of len can change after we read it and before we set it!
- Solution
 - Lock the variable before doing “test and set”



- Every Java object has a wait/notify queue
 - Similar to the way every Java object has a lock
 - Used to get threads to cooperate with or signal each other
- The queue is like the *len* variable in the previous example!
 - i.e. we MUST have a lock on the object before we can touch it's queue
 - Implies that wait() and notify can only be called inside a synchronized method or a synchronized(obj) {...} block
 - Must synchronize on the object whose queue is being used!



- obj.wait();
 - Send to any object
 - Calling thread waits (blocks) on the object's queue
 - Efficient blocking
 - Must first have that objects lock!
 - Waiting thread releases that objects lock
 - Does not release any other locks it holds!
 - Sending an interrupt() to the waiting thread will result in popping out of its wait
 - Actually this will result in a InterruptedException which would need to be caught
 - We will see this in an example later



- obj.notify(); obj.notifyAll();
 - Send to any object
 - Notifies a waiter (thread) on that objects queue if there is one
 - Sender must have the objects lock
 - A random waiting thread will get woken up from its wait()
 - Not necessarily FIFO
 - Not right away
 - Waiter will re-acquire the lock before resuming operation



- Dropped notify()
 - If there are no waiting threads on the objects queue, the notify() does nothing
 - wait()/notify() **do not count up and down**
 - That requires a semaphore – see handout
- notifyAll()
 - Notifies all waiting threads on the queue
 - Tricky to know when to call notify()
 - Most common approach is to always call notifyAll()
 - Only one thread will be able to acquire the lock
 - Not too expensive



Monitor Exception

- `java.lang.IllegalMonitorStateException`: current thread not owner
 - This is the exception thrown if a thread tries a `wait/notify` on a object without first holding its lock!
 - You will get these while writing your code!
 - Make sure you are synchronizing on the correct object before calling `wait` or `notify`!



While (cond) wait() idiom

- When the waiting thread is woken up from the `wait` it holds the lock
 - But the condition it was waiting for may not be true any more!
 - It may have become false again in between when the `notify` happened and when the `wait/return` happened
 - Necessary to check the condition again before proceeding further
- Idiom


```
try {
    while (<condition>) wait();
} catch (InterruptedException e) {}
```



Wait/notify example

- Producer/Consumer problem with `wait/notify`
 - "len" represents the number of elements in some imaginary array
 - `add()` adds an element to the end of the array. `Add()` never blocks
 - We assume there's enough space in the array.
 - `remove()` removes an element, but can only finish if there is an element to be removed.
 - If there is no element, `remove()` waits for one to be available.



Wait/notify example

- Strategy:
 - The `AddRemove` object is the common object between the threads
 - they use its lock and its `wait/notify` queue.
 - `add()` does a `notify()` when it adds an element
 - `remove()` does a `wait()` if there are no elements
 - Eventually, an `add()` thread will put an element in and do a `notify()`
 - Each adder adds 10 times, and each remover removes 10 times, so it balances in the end.



Wait/Notify example code

- Code walk through
 - In emacs...



Dropped notify() problem...

- `Notify()` does not count the number of notifies!
 - It is instantaneous
 - If there are waiters waiting they will be signaled
 - If a waiter comes after the `notify`, it is not signaled
- `wait/notify()` is simpler than a semaphore
 - Semaphores count
 - Classic CS locking construct
 - Possible to build semaphore using `wait/notify`



- Code walkthrough
 - In emacs...



- Problem: Swing vs. Threads
 - Modifying the GUI state while it is being drawn
 - Typical reader/writer conflict problem
 - Example
 - paintComponent() while another thread changes the component geometry
 - Send mouseMoved() notification to an object, but another thread deletes the object!



- Swing Thread: a.k.a One Big Lock!
 - One official designated “Swing thread”
 - Does all Swing/GUI notifications using the Swing thread, one at a time
 - paintComponent() – always on Swing Thread
 - All notifications: action events, mouse events – sent on the Swing Thread
 - System keeps a queue of “Swing jobs”
 - When the Swing Thread is done with its current job it moves on to the next one
 - Only the Swing Thread is allowed to edit the state of the GUI
 - Since the Swing thread is the only one allowed to touch the Swing state there is in effect a big lock over all the Swing State



- On the swing thread – edit ok
 - Allowed to edit the Swing state when you are on the Swing Thread
 - Container.add(), setPreferredSize(), setLayout()
- Don't hog the Swing Thread
 - Do not to time-consuming operations on the Swing Thread
 - If you hold the Swing Thread, no Swing/GUI processing will happen till you release it!
 - Fork off a worker thread to do a time-consuming operation
- Not on the Swing Thread – no edit
 - A thread which is not the swing thread may not send messages that edit the Swing state
 - Use invokeLater() to run code on the swing thread
 - Repaint() is an exception – since it only schedules a call to paintComponent() which is called by the Swing Thread
 - Another exception is modifying state before the component has been made visible
 - For example in a constructor



- In your notifications (paintComponent(), actionPerformed()) you are on the Swing Thread
 - Feel free to send Swing messages
- There is only one Swing Thread
 - When you have it, no other Swing activity is happening
 - Do not hog the Swing Thread



- Built in utility method to allow you to “post” some code to the Swing Thread to run later
 - Uses Runnable interface
 - public void run()
 - SwingUtilities.invokeLater(Runnable)
 - Queue up the given runnable
 - Will execute when the Swing Thread gets to it
 - SwingUtilities.invokeAndWait(Runnable)
 - Same as above, but also block current thread till the runnable has completed

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SwingUtilities Client Example

```

class MyFrame extends JFrame {
    private JLabel label;
    // Typical GUI code down here creates and starts the worker
    public MyFrame() {
        // standard Frame ctor stuff, create buttons...
        button.addActionListener( new ActionListener() {
            public void actionPerformed(ActionEvent e) {
                Worker worker = new Worker();
                worker.start();
            }
        });
        ...
    }
}

```

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SwingUtilities Client Example

```

class Worker extends Thread {
    public void run() {
        // The worker does some big computation
        final String answer = <something>;
        // We want to call setText() to send the answer to the GUI.
        // We must go through SwingUtilities.invokeLater()
        SwingUtilities.invokeLater(
            new Runnable() { // create a runnable on the fly
                public void run() {
                    label.setText(answer);
                }
            }
        );
    }
}


```

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SwingThread Demo

- Demo and code walkthrough
 - In emacs...



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Threading Conclusions

- Java uses an OOP Concurrency style
 - Objects store state
 - Getters and setters are synchronized
 - Intuitive extension to how threading is handled
 - Not just a translation from C/C++
- Compile Time "Structured" style
 - Lock/unlock structure is specified at compile time
 - synchronized(obj) {...}
 - Impossible to write code where lock/unlock don't balance
 - Much better than lock() and unlock() constructs in other languages
- Java uses "monitor" style locking
 - Not as flexible, but easier and less error prone

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When to use Threading?

- Hardware
 - To take advantage of increasingly parallel hardware
- GUI
 - To keep the GUI responsive
- Networking
 - Use thread to support multiple connections
 - Speed up by pipelining slow operations

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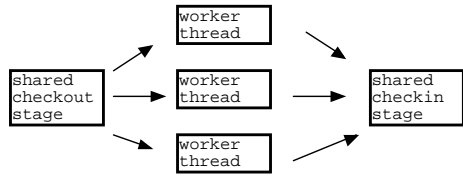
In general...

- Single Threaded is easier!
 - There are cases when this is the best use of your time
- Design for concurrency
 - By default, do not put much effort in to making your class support concurrency
 - Should only be deliberately added when it makes sense
 - It is not trivial to support concurrency
 - Performance tradeoff
 - Complexity tradeoff

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Typical Good Design – Checkin/Checkout



Summary

- Today
 - Thread Interruption
 - Cooperation
 - Wait/notify
 - Swing/GUI Threading
 - SwingThread Demo
 - Threading conclusions
- Assigned Work Reminder
 - HW 3a: ThreadBank
 - HW 3b: LinkTester
 - Both due before midnight on Wednesday, August 6th, 2003
 - Do HW3a this week!!